# Parallel Breadth-First Search Using OpenMP

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# **Top-down Approach**

**Implementation Details** 

We use most of the same implementation detail as the sequential top-down solution provided: two vertex\_set (frontier and new\_frontier). We parallelize the BFS by adding #pragma omp parallel for, and made the frontier check atomic, using \_\_sync\_bool\_compare\_and\_swap. However, since we don't even need to try to swap if the entry is visited before, we do a test before calling \_\_sync\_bool\_compare\_and\_swap to reduce synchronization. Also, if the swap is successful, we use a \_\_sync\_add\_and\_fetch to atomically add the counter in the new frontier. Also, we passed a few constants and pointers into top\_down\_step function like graph -> num\_nodes and graph -> outgoing\_starts, so that the threads don't need to access memory to fetch them every time.

## **Optimization Process**

First we implemented details above except doing the test before \_\_sync\_bool\_compare\_and\_ swap. The performance of our solution is not quite stable, and the performance can go higher than 125% of reference solution sometimes. Then we went to office hours and realized that synchronization could take a lot of costs, so we started to think about what synchronization steps can be "eliminated". After a while we came up with the idea that we don't really need to swap the entry in distance array if it's already visited, and thus add a test before calling the atomic function.

## **Performance analysis**

- Where is the synchronization in your solution? Do you do anything to limit the overhead of synchronization?

There are two places that synchronization happens: to (atomically) compare and swap in the distance array, and to (atomically) add the counter in new\_frontier. We did a test before doing the compare and swap in order to reduce synchronization cost.

Why do you think your code is unable to achieve perfect speedup?

The workload should be relatively balanced for top-down approach because frontiers are evenly divided among all processors, and every frontier should contain the same amount of work on average. However, there are considerable communication/synchronization in the top down approach (two synchronization steps). There isn't much artificial data movement in top down step. - When you run your code on Blacklight with more than 16 threads, do you see a noticeable drop off in performance? Why do you think this might be the case?

We didn't see a noticeable drop off in performance of top-down approach when the thread counts go over 16 on Blacklight.

## **Bottom-up Approach**

#### **Implementation Details**

We use two boolean arrays (frontier and new\_frontier) to keep track of frontiers and new frontiers. There is no synchronization in this approach. For every vertex, we can check the distance array to see if it has been visited before, and if not, we can go through all its incoming neighbors, and check frontier array to see if any neighbor was visited before. If it is, then it must be a new frontier, so we set the respective entry in new\_frontier to true, and update this node's distance. We used "guided" scheduling for the approach. In definition, guided is pretty much like dynamic, but at the start of the program every thread will be assigned more work than designated, and gradually every thread will be assigned less and less work, with a minimum of the designated number. We think this approach will decrease the scheduling overhead for large graphs. Also, we passed a few constants and pointers into bottom\_up\_step function like graph->num\_nodes and graph-> incoming\_starts, so that the threads don't need to access memory to fetch them every time.

#### **Optimization Process**

Reading the spec, we first realize that doing linear check to see if a vertex was visited before, or if its incoming neighbors were visited before is probably going to be too slow. Thus, the first solution we came up with is one that uses distances array to check both conditions in O(1). Also there were no synchronizations in the solution, because every node only write to its own entry in distances array. Surprisingly the solution was not fast enough. Confused, we went to office hour. After the TA read several similar solutions that look alike, he concluded that reading and writing one array for all threads may suffer from false sharing, and suggested us to add more data structures to separate read and write. Thus we came up the approach to use frontier (read-only in a round) and new\_frontier (write-only in a round), and use boolean array for both in order to increase the number of items that can be held in a cache line. After implementing the optimization above, the performance of our bottom-up approach increased a lot, but still can be unsatisfactory for rmat graphs. Thus, we made an extra optimization to pass constants and pointers to bottom\_up\_step so that they don't need to be dereferenced every time to be used, and the performance of rmat graphs became better.

**Performance analysis** 

- Where is the synchronization in your solution? Do you do anything to limit the overhead of synchronization?

There is no explicit synchronization in the solution (no atomic function calls/critical sections). However, cache coherence problem can happen sometimes, because many professors can potentially read from and write to the same cache line. We use several redundant arrays to separate read and write in order to reduce false sharing.

Why do you think your code is unable to achieve perfect speedup?

The workload should be relatively balanced for in our approach, because we used dynamic scheduling. However, there are some implicit synchronization/communication costs (false sharing). These costs can also be attributed to data movement because when cache are invalidated, data move from one cache to main memory/another cache.

- When you run your code on Blacklight with more than 16 threads, do you see a noticeable drop off in performance? Why do you think this might be the case?

Yes we did. Since the main cost of bottom-up approach is from cache coherence, adding the number of cores could hurt speedup.

# Hybrid Approach

**Implementation Details** 

The intuition for hybrid is that top-down is fast when there are few frontiers (because it's likely that there are not many new frontiers, so the bottom-up's approach will waste time traversing all vertices), and bottom-up is fast when there are many frontiers (because top-down will have many new frontiers overlap, increasing synchronization overhead and repeated checks).

We use the same structure and algorithm for hybrid as our parallel bottom-up and top-down solution, except that we keep track of bottom-up's frontier array when we are doing top-down. Once we switch to bottom-up when there are many frontiers, we never switch back again because switching between data structures are expensive. Our policy of switching is that if the ratio of the number of frontiers to the total number of vertices is greater than 0.02, we do the switching. This "magic number" is obtained by experiment.

## **Optimization Process**

First our switching policy is that if the number of frontiers exceed a certain number, then we do the switching. However, this solution is not scalable because graph size (and thus frontier size) could differ greatly. Therefore, we changed our policy to ratio, and after several runs, we decided that 0.02 is a good ratio, favoring large graphs.

Also, we keep track of both approaches' data structures so we can switch back and forth. However, this method proved to be too slow because keeping track of top-down approach's data structures requires synchronization. Then we tried converting one type of data structure to the other when we are about to do the switching, and it still proved to be too slow. Finally, we realized that when the frontier gets large, there shouldn't be too many iterations before the graph is completed processed, so it doesn't hurt too much to keep using bottom-up and never switch back. This approach proved to be sufficiently fast.

# **Performance analysis**

- Where is the synchronization in your solution? Do you do anything to limit the overhead of synchronization?

The only explicit synchronization is the synchronization for top-down approach in our solution. We didn't do any extra step compared to top-down approach. However, we decided not to convert top-down's data structure to that of bottom-up when we want to do the switching, but instead keep track of it when we are doing top-down to reduce the synchronization step needed to convert. Also we decide not to switch back to top-down approach after we take bottom-up because either keeping track of top-down approach's data structures or converting bottom-up approach's data structures to that of top-down could add synchronization costs.

- Why do you think your code is unable to achieve perfect speedup?

The workload should be balanced according to the analysis in the previous two approaches. There are explicit synchronization costs in top-down approach, and there are implicit cache communication costs in both approaches. Also, there are extra data movement in top-down approach because it needs to keep track of bottom-up approach's data structure.

- When you run your code on Blacklight with more than 16 threads, do you see a noticeable drop off in performance? Why do you think this might be the case?

We saw some drop off in performance of hybrid approach when the thread counts go over 16 on Blacklight, though it not as noticeable as the bottom-up approach. We think the reason is that the drop off is counterbalanced by the top-down approach

# **Tables:**

Runtime on ghc39.ghc.andrew.cmu.edu (rmat\_32m.graph):

Timing S	Summary					
Threads	Top Down		Botto	om Up	Hybi	rid
1:	6.0032 (1.	0000x)	6.9139	(1.0000x)	3.4972	(1.0000x)
2:	4.0748 (1.	4732x)	3.8161	(1.8118x)	1.9381	(1.8045x)
4:	2.9347 (2.	0456x)	2.2236	(3.1093x)	1.1487	(3.0445x)
6:	2.6054 (2.	3041x)	1.8216	(3.7956x)	0.9440	(3.7046x)
8:	2.3794 (2.	5230x)	1.7619	(3.9242x)	0.8942	(3.9112x)
12:	2.0224 (2.	9684x)	1.6827	(4.1089x)	0.8406	(4.1603x)
Referenc	e Summary					
Threads	Top Down		Botto	om Up	Hybi	rid
1:	6.4014 (1.	0000x)	7.8318	(1.0000x)	3.8490	(1.0000x)
2:	4.2107 (1.	5203x)	4.3779	(1.7889x)	2.2216	(1.7325x)
4:	3.3601 (1.	9051x)	2.5808	(3.0346x)	1.4140	(2.7221x)
6:	3.2257 (1.	9845x)	2.1347	(3.6688x)	1.2215	(3.1511x)
8:	2.8480 (2.	2477x)	2.0497	(3.8209x)	1.2122	(3.1752x)
12:	2.3861 (2.	6828x)	1.9423	(4.0322x)	1.1111	(3.4642x)
For grac Correctr	ling refere	ence (base	d on exec	cution times	)	
Timina:						
Threads	Top Down	Bot	tom Up	Hvbrid		
1:	93.78p	88.28p	90.86	0		
2:	96.77p	87.17p	87.24	0		
4:	87.34p	86.16p	81.24	0		
6:	80.77p	85.33p	77.29	0		
8:	83.55p	85.96p	73.76	0		
12:	84.76p	86.63p	75.66	0		

Runtime on ghc39.ghc.andrew.cmu.edu (random\_50m.graph):

Timina	Summary					
Threads	Top Down		Bottom	Up	Hybri	d
1 :	14,6249 (1	0000x)	34,1084	(1.0000x)	5 6928	(1.0000x)
2.	9 6571 (1	5144x)	21 0551	(1 6200x)	3 6465	(1.5611x)
<u>z</u> .	6 9414 (2	1069x1	12 8242	2 6597x)	2 8747	(1 9803x)
6.	5 9244 (2	4686×)	15 9294	(2.1412)	2 3804	(2 3015v)
g .	5 5000 (2	6590x)	17 0268	2 00322)	2 2875	(2.4886v)
12:	6.0930 (2.	4003x)	10.3422	(3.2980x)	2.3312	(2.4420x)
Referen	ce Summarv					
Threads	Top Down		Bottom	Up	Hybri	d
1:	15.7651 (1	.0000x)	36,7491	(1.0000x)	8.7110	(1.0000x)
2	10.5329 (1	.4967x)	25,2021	(1,4582x)	4.7253	(1.8435x)
4	7.5428 (2.	0901x)	17,9428	(2.0481x)	3,2726	(2.6618x)
6 :	6.9795 (2.	2588x)	16.3647	2.2456x)	3,1684	(2.7493x)
8:	6.2725 (2.	5134x)	17.0879	2.1506x)	3.2396	(2.6889x)
12:	6.3549 (2.	4808x)	16.3004	(2.2545x)	2.8149	(3.0946x)
For grad	dina refere	nce (base	d on execut	tion times)		
Correct	noss.					
correcti	1633.					
Timing:						
Threads	Top Down	Bot	tom Up	Hybrid		
1:	92.77p	92.81p	65.35p			
2:	91.68p	83.54p	77.17p			
4:	92.03p	71.47p	87.84p			
6:	84.88p	97.34p	75.13p			
8:	87.68p	99.64p	70.61p			
12:	95.88p	63.45p	82.82p			

Top Down Bottom Up Hybrid	
Bottom Up 11.6702	Hybrid 6.9809
Bottom Up 14.0961	Hybrid 7.2650
2	
Top Down Bottom Up Hybrid	
Bottom Up 10.9551	Hybrid 4.4316
Bottom Up 8.1047	Hybrid 4.8431
	Top Down Bottom Up Hybrid Bottom Up 11.6702 Bottom Up 14.0961 2 P Top Down Bottom Up Hybrid Bottom Up Hybrid Bottom Up Hybrid Bottom Up 10.9551 Bottom Up 8.1047

Runtime on Blacklight (rmat\_32m.graph):

Max system threads = 4 Running with 4 threads		
Loading graph		
Graph stats: Edges: 199491925 Nodes: 33554432 Running with 4 threads Testing Correctness of T Testing Correctness of F Testing Correctness of F	°op Down Bottom Up Hybrid	
Timing Summary Threads Top Down 4: 11.4126	Bottom Up 8.6600	Hybrid 2.6278
Reference Summary Threads Top Down 4: 11.2683	Bottom Up 4.0421	Hybrid 2.7536
Max system threads = 8 Running with8 threads		
Loading graph		
Graph stats: Edges: 199491925 Nodes: 33554432 Running with 8 threads Testing Correctness of T	Fop Down	
Testing Correctness of F	Iybrid	
Timing Summary Threads Top Down 8: 8.3346	Bottom Up 7.5327	Hybrid 1.6313
Reference Summary Threads Top Down 8: 11.2405	Bottom Up 2.1812	Hybrid 1.7437

Max system threads = 1 Running with 16 threads	6	
Loading graph		
Graph stats: Edges: 199491925 Nodes: 33554432 Running with 16 threads Testing Correctness of T Testing Correctness of F Testing Correctness of F	s Fop Down 3ottom Up 4ybrid	
Timing Summary Threads Top Down 16: 10.6528	Bottom Up 6.5767	Hybrid 1.9410
Reference Summary Threads Top Down 16: 11.3013	Bottom Up 1.4047	Hybrid 1.5714
Max system threads = 3 Running with 32 threads	2 3	
Loading graph		
Graph stats: Edges: 199491925 Nodes: 33554432 Running with 32 threads Testing Correctness of T Testing Correctness of F Testing Correctness of F	s Fop Down Bottom Up Hybrid	
Timing Summary Threads Top Down 32: 24.1170	Bottom Up 21.5944	Hybrid 3.9215
Reference Summary Threads Top Down 32: 24.6346	Bottom Up 4.6247	Hybrid 4.4369

Max system threads = 1 Running with 1 threads	l	
Loading graph		
Graph stats: Edges: 499999944 Nodes: 50000000 Running with 1 threads Testing Correctness of Testing Correctness of Testing Correctness of	Top Down Bottom Up Hybrid	
Timing Summary Threads Top Down 1: 46.3898	Bottom Up 55.7806	Hybrid 11.2587
Reference Summary Threads Top Down 1: 44.9548	Bottom Up 68.4181	Hybrid 15.2585
Max system threads = 2 Running with 2 threads	2	
Loading graph		
Graph stats: Edges: 499999944 Nodes: 50000000 Running with 2 threads Testing Correctness of Testing Correctness of Testing Correctness of	Top Down Bottom Up Hybrid	
Timing Summary Threads Top Down 2: 36.3385	Bottom Up 35.1577	Hybrid 7.7612
Reference Summary Threads Top Down 2: 36.2882	Bottom Up 36.0831	Hybrid 10.1436

Runtime on Blacklight (random\_50m.graph):

Max system threads = 4 Running with 4 threads		
Loading graph		
Graph stats: Edges: 499999944 Nodes: 5000000 Running with 4 threads Testing Correctness of T Testing Correctness of B Testing Correctness of H	op Down ottom Up lybrid	
Timing Summary Threads Top Down 4: 21.6723	Bottom Up 24.1341	Hybrid 4.8610
Reference Summary Threads Top Down 4: 21.5067	Bottom Up 18.7704	Hybrid 5.6025
Max system threads = 8 Running with 8 threads		
Loading graph		
Graph stats: Edges: 499999944 Nodes: 5000000 Running with 8 threads Testing Correctness of T Testing Correctness of B Testing Correctness of H	op Down ottom Up (vbrid	
Timing Summary		
Threads Top Down 8: 13.4517	Bottom Up 15.6372	Hybrid 2.9507
Reference Summary Threads Top Down 8: 17.6449	Bottom Up 10.1982	Hybrid 3.9919

Max system threads = 1 Running with 16 thread	6 s	
Loading graph		
Graph stats: Edges: 499999944 Nodes: 5000000 Running with 16 threads Testing Correctness of T Testing Correctness of I Testing Correctness of I	s Fop Down Bottom Up Hybrid	
Timing Summary Threads Top Down 16: 19.3297	Bottom Up 17.6897	Hybrid 4.1888
Reference Summary Threads Top Down 16: 20.6625	Bottom Up 7.5523	Hybrid 3.8894
Max system threads = 3 Running with 32 threads	2 s	
Loading graph		
Graph stats: Edges: 499999944 Nodes: 5000000 Running with 32 thread Testing Correctness of T Testing Correctness of I Testing Correctness of I	s Fop Down Bottom Up Hybrid	
Timing Summary Threads Top Down 32: 46.4034	Bottom Up 31.5984	Hybrid 10.0992
Reference Summary Threads Top Down 32: 56.1906	Bottom Up 39.2686	Hybrid 13.8575

Runtime on unix.andrew.cmu.edu (random\_50m.graph):

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Timina	Summary					
Threads	Top Down		Bottom	Up	Hvb	rid
1 :	15.8681 (1	.0000x)	25.8125	(1.0000x)	4.48	28 (1.0000x)
2:	13,4549 (1	.1794x)	17.0456	(1.5143x)	3.40	20 (1.3177x)
4.	9 1159 (1	7407×1	9 1072 (2	8343x1	2 0885	(2 1464x)
8.	6 9225 (2	2922×1	5 2636 (4	9040x)	1 5145	(2 9599x)
16.	5 8690 (2	7037x)	3 3598 (7	6828x)	1 2651	(3.5434x)
32:	4.6955 (3.	3794x)	3.3549 (7	(.6939x)	1.0765	(4.1641x)
Reference	ce Summary					
Threads	Top Down		Bottom	Up	Hvb	rid
1:	16.3994 (1	.0000x)	27.4077	(1.0000x)	5.70	(1.0000x)
2:	14.3532 (1	.1426x)	17.7396	(1.5450x)	4.394	47 (1.2973x)
4:	10.7889 (1	.5200x)	10.2806	(2.6660x)	2.84	10 (2.0046x)
8:	9.2487 (1.	7732x)	6.9772 (3	3.9282x)	2.1605	(2.6388x)
16:	8.2343 (1.	9916x)	5.4853 (4	.9966x)	1.8675	(3.0527x)
32:	6.2013 (2.	6445x)	5.2117 (5	.2588x)	1.5797	(3.6091x)
For grad	ding refere	ence (base	d on execut	ion times)		
Correcti	ness:					
Tinina						
Timing:	THE OWNER STREET	Det	12/12/22/11 10:22:10	ELECTRON CONTRACTOR		
inreads	OF 765	04 10 BOL		пуртта		
1:	90.70p	94.18p	78.03p			
2.	93.74p	96.09p	77.41p			
4:	84.49p	88.59p	73.44p			
8:	74.85p	75.44p	70.10p			
10:	71.2/p	61.25p	67.74p			
32:	/5./2p	64.3/p	68.15p			